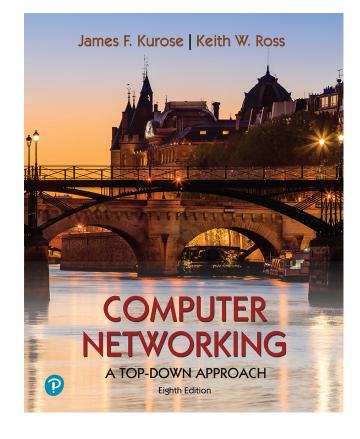
Chapter 3 Transport Layer



Computer Networking: A Top-Down Approach

8th edition Jim Kurose, Keith Ross Pearson, 2020

Transport layer: roadmap

- Transport-layer services
- Multiplexing and demultiplexing
- Connectionless transport: UDP
- Principles of reliable data transfer
- Connection-oriented transport: TCP
- Principles of congestion control
- TCP congestion control
- Evolution of transport-layer functionality



Evolving transport-layer functionality

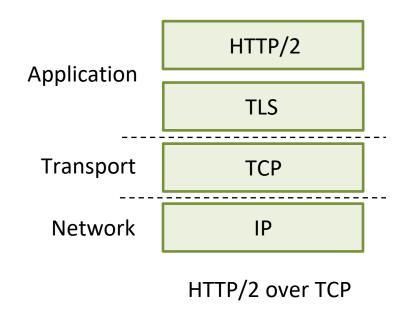
- TCP, UDP: principal transport protocols for 40 years
- different "flavors" of TCP developed, for specific scenarios:

Scenario	Challenges
Long, fat pipes (large data	Many packets "in flight"; loss shuts down
transfers)	pipeline
Wireless networks	Loss due to noisy wireless links, mobility;
	TCP treat this as congestion loss
Long-delay links	Extremely long RTTs
Data center networks	Latency sensitive
Background traffic flows	Low priority, "background" TCP flows

- moving transport—layer functions to application layer, on top of UDP
 - HTTP/3: QUIC

QUIC: Quick UDP Internet Connections

- application-layer protocol, on top of UDP
 - increase performance of HTTP
 - deployed on many Google servers, apps (Chrome, mobile YouTube app)

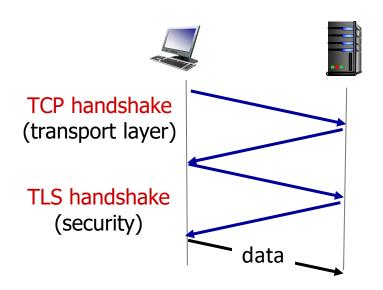


QUIC: Quick UDP Internet Connections

adopts approaches we've studied in this chapter for connection establishment, error control, congestion control

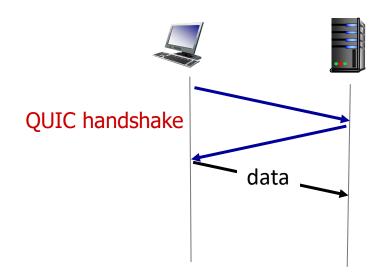
- error and congestion control: "Readers familiar with TCP's loss detection and congestion control will find algorithms here that parallel well-known TCP ones." [from QUIC specification]
- connection establishment: reliability, congestion control, authentication, encryption, state established in one RTT
- multiple application-level "streams" multiplexed over single QUIC connection
 - separate reliable data transfer, security
 - common congestion control

QUIC: Connection establishment



TCP (reliability, congestion control state) + TLS (authentication, crypto state)

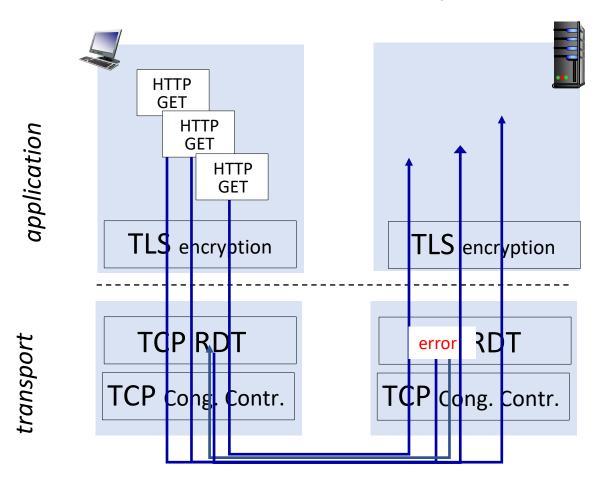
2 serial handshakes



QUIC: reliability, congestion control, authentication, crypto state

1 handshake

QUIC: streams: parallelism, no HOL blocking



(a) HTTP 1.1

Chapter 3: summary

- principles behind transport layer services:
 - multiplexing, demultiplexing
 - reliable data transfer
 - flow control
 - congestion control
- instantiation, implementation in the Internet
 - UDP
 - TCP

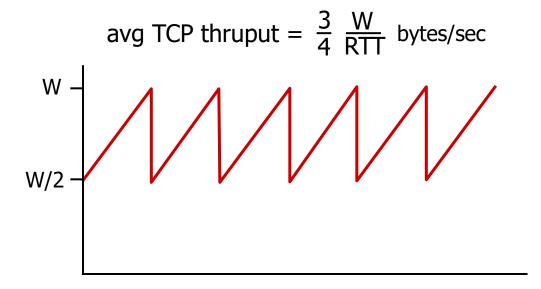
Up next:

- leaving the network "edge" (application, transport layers)
- into the network "core"
- two network-layer chapters:
 - data plane
 - control plane

Additional Chapter 3 slides

TCP throughput

- avg. TCP thruput as function of window size, RTT?
 - ignore slow start, assume there is always data to send
- W: window size (measured in bytes) where loss occurs
 - avg. window size (# in-flight bytes) is ¾ W
 - avg. thruput is 3/4W per RTT



TCP over "long, fat pipes"

- example: 1500 byte segments, 100ms RTT, want 10 Gbps throughput
- requires W = 83,333 in-flight segments
- throughput in terms of segment loss probability, L [Mathis 1997]:

TCP throughput =
$$\frac{1.22 \cdot MSS}{RTT \sqrt{L}}$$

- → to achieve 10 Gbps throughput, need a loss rate of L = $2 \cdot 10^{-10} a$ very small loss rate!
- versions of TCP for long, high-speed scenarios