

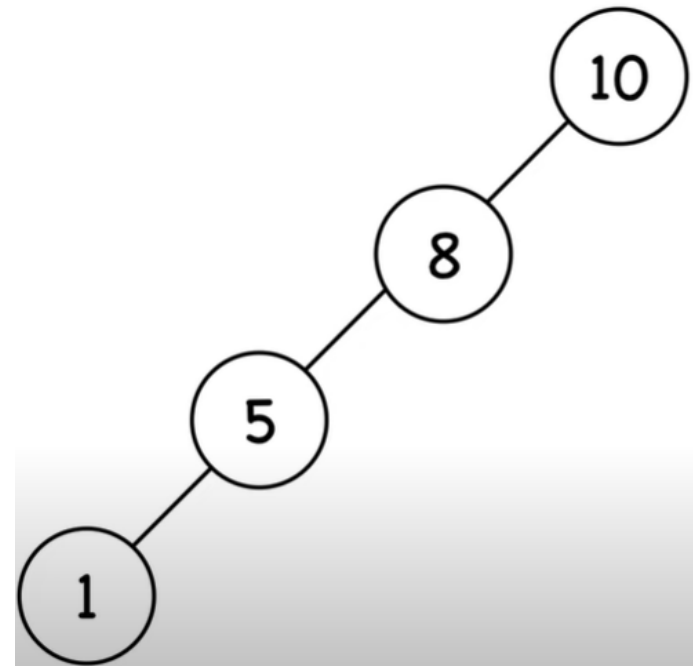
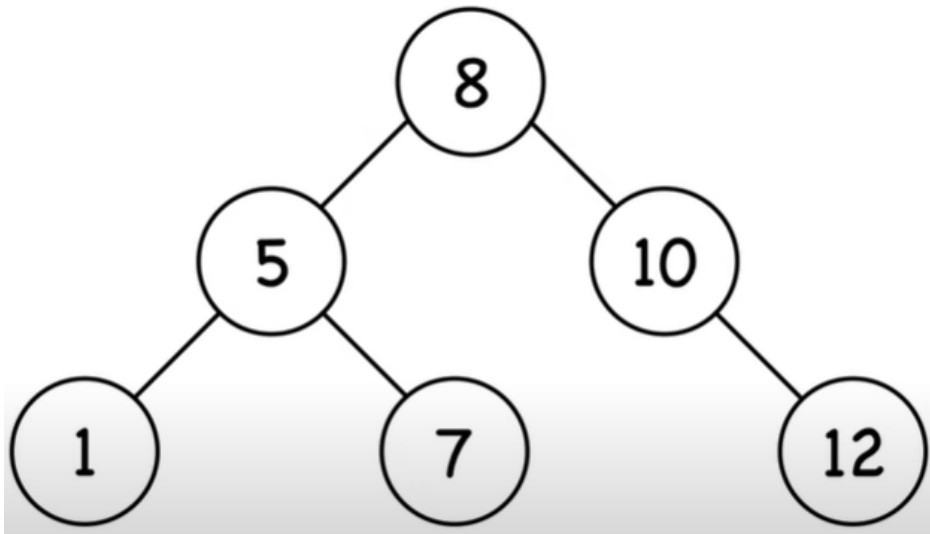
Lecture 9

Red-Black Trees

Department of Computer Science
Hofstra University

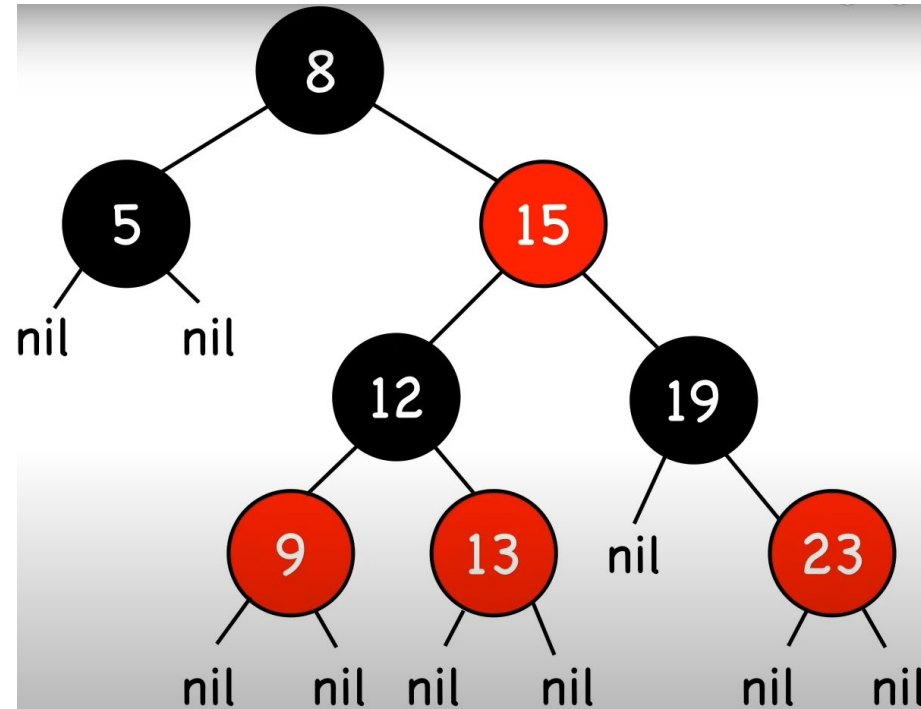
Binary Search Trees

- Ordered, or sorted, binary trees.
- Each node can have 2 subtrees.
- Items to the left of a given node are smaller.
- Items to the right of a given node are larger.
- Balanced search trees have guaranteed height of $O(\log n)$ for n items
 - Red-Black Tree is a type of balanced search tree



Red-Black Tree

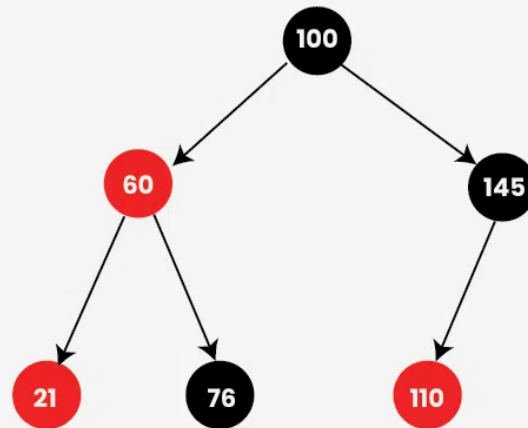
- 1. Node Color: A node is either red or black.
- 2. Root Property: The root and leaves (NIL) are black.
- 3. Red Property: If a node is red, then its children are black.
- 4. Black Property: All paths from a node to its NIL descendants contain the same number of black nodes.
 - Path length excludes root node itself, so here each path contains 1 black node



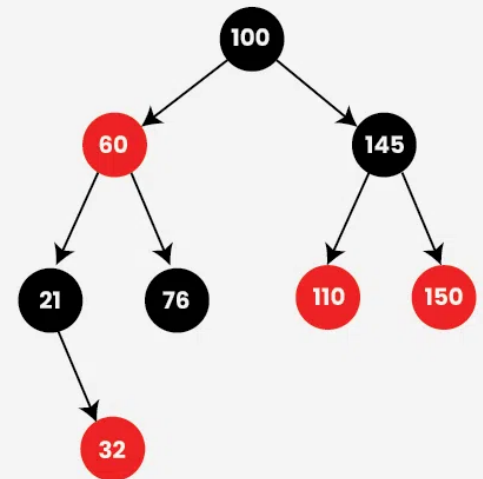
Example

- Tree on the left:
Incorrect Red Black Tree.
 - Two red nodes are adjacent to each other.
 - One of the paths to a leaf node has zero black nodes, whereas the other two paths contain 1 black node each.

Example of Red-black Tree



A incorrect Red-black Tree

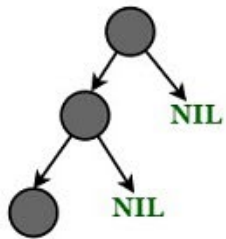


A correct Red-black Tree

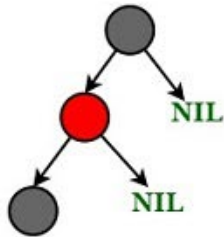
Red-Black tree ensures balancing

- A chain of 3 nodes is not possible in a Red-Black tree

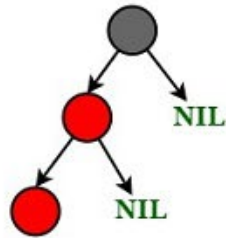
Following are NOT possible 3-noded Red-Black Trees



Violates
Property 4



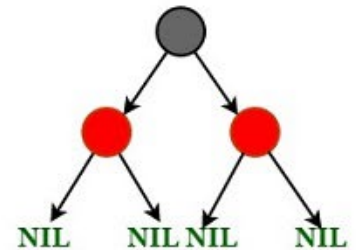
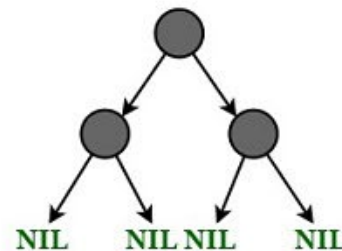
Violates
Property 4



Violates
Property 3



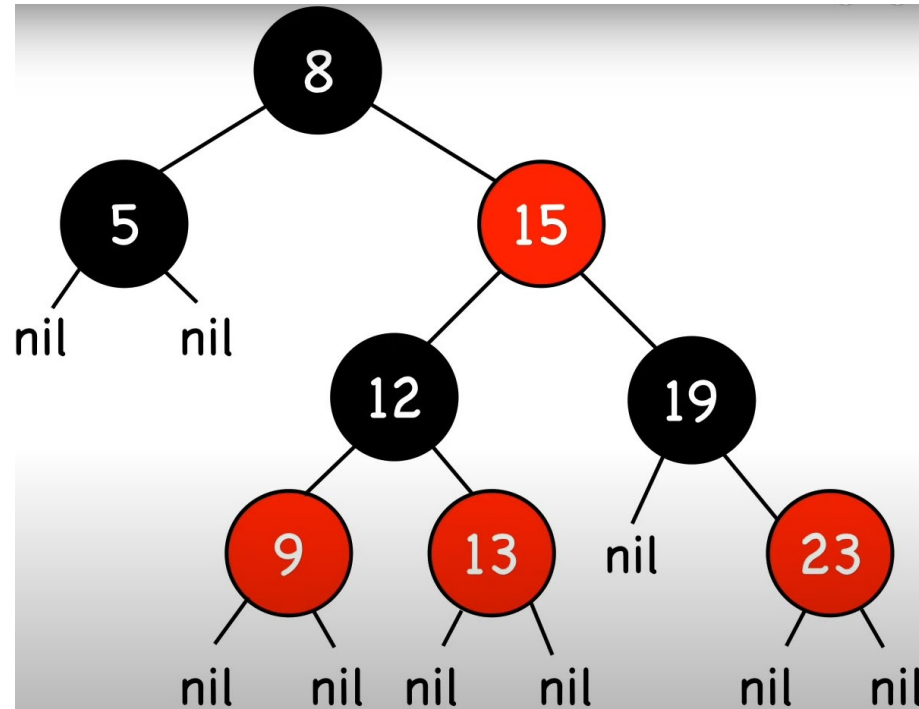
Following are possible Red-Black Trees with 3 nodes



All Possible Structure of a 3-noded Red-Black Tree

Additional Properties

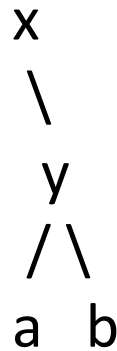
- Balanced search tree: the longest path (root to farthest NIL) is no more than twice the length of the shortest path (root to nearest NIL).
 - Shortest path: all black nodes (=2)
 - Longest path: alternating red and black (=4)
- Operations: search, insert, remove, each with time complexity $O(\log(n))$.
 - Insert and remove may result in violation of red-black tree properties, use rotations to fix it



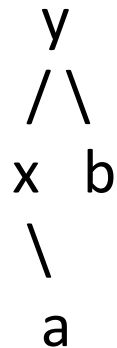
Rotations

- Alters the structure of a tree by rearranging subtrees
- Goal is to decrease the height of the tree to maximum height of $O(\log n)$
 - Larger subtrees up, smaller subtrees down
- Does not affect the order of elements
- Time complexity $O(1)$

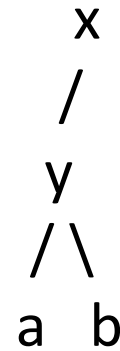
Before Rotation:



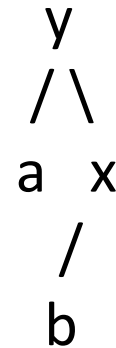
After Left Rotation:



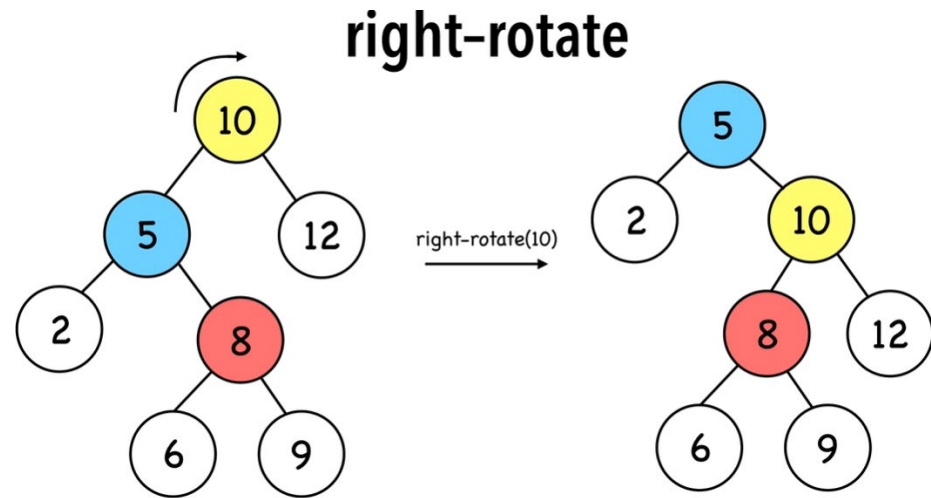
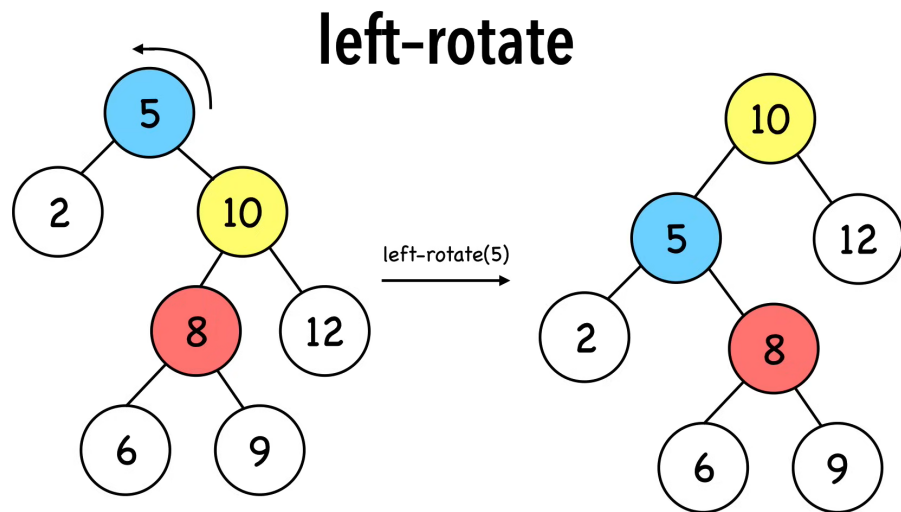
Before Rotation:



After Right Rotation:



Rotations Examples

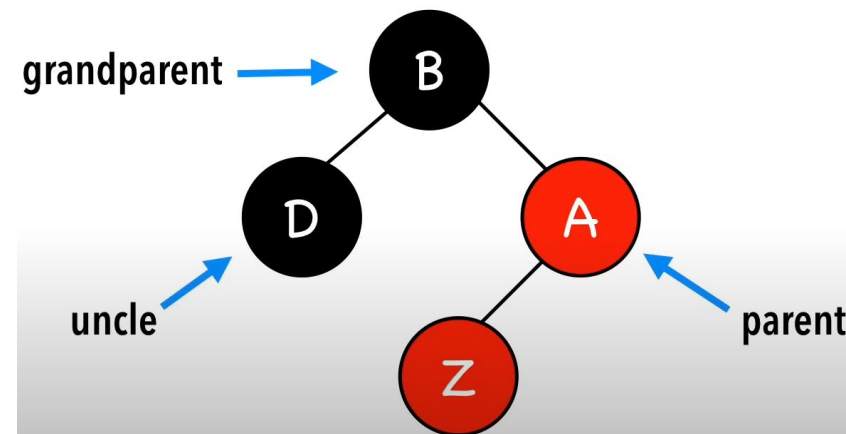


Insertion

- Inserting a new node in a Red-Black Tree involves a two-step process: performing a standard binary search tree (BST) insertion, followed by fixing any violations of Red-Black properties.
- **Insertion Steps**
 1. **BST Insert:** Insert the new node into BST and color it red.
 2. **Fix Violations:**
 2. If the parent of the new node is **black**, no properties are violated.
 3. If the parent is **red**, the tree might violate the Red Property, requiring fixes.

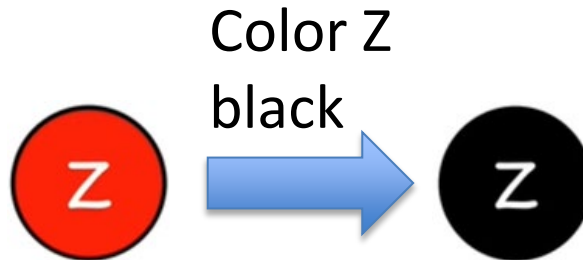
Insertions

- Step 1. Insert Z and color it **red**
- Step 2. Recolor and rotate nodes to fix violations
- 4 scenarios after inserting node Z
- Case 0. Z = root
 - Color Z black
- Case 1. Z.uncle = **red**
 - Recolor Z's parents and grandparent
- Case 2. Z.uncle = black (triangle)
 - Rotate Z.parent, turns into Case 3
- Case 3. Z.uncle = black (line)
 - Rotate Z.grandparent & Recolor Z's parents and grandparent



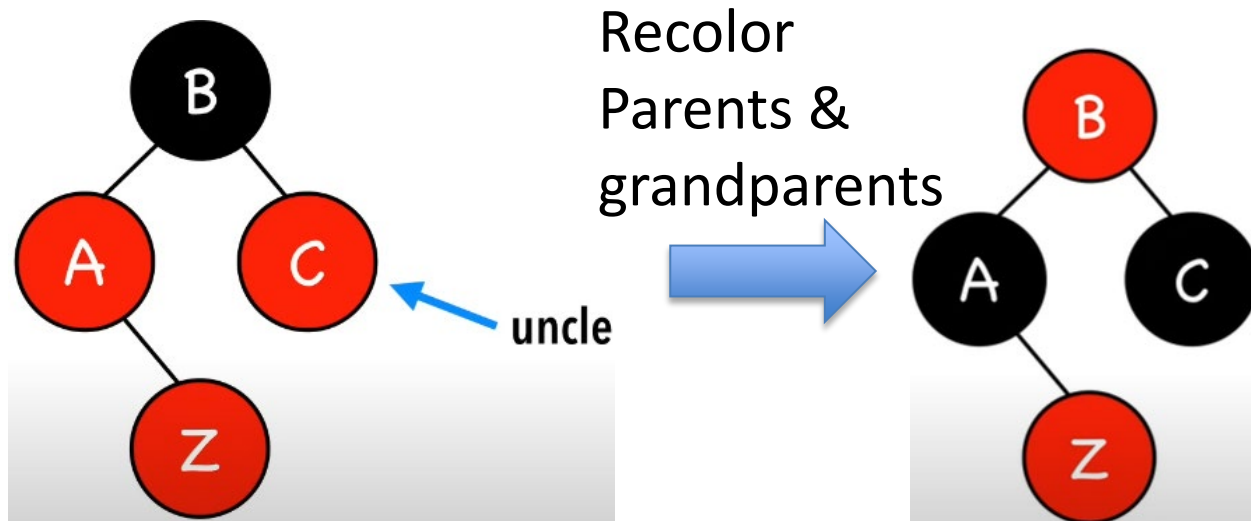
Case 0. $Z = \text{root}$

- Color Z black



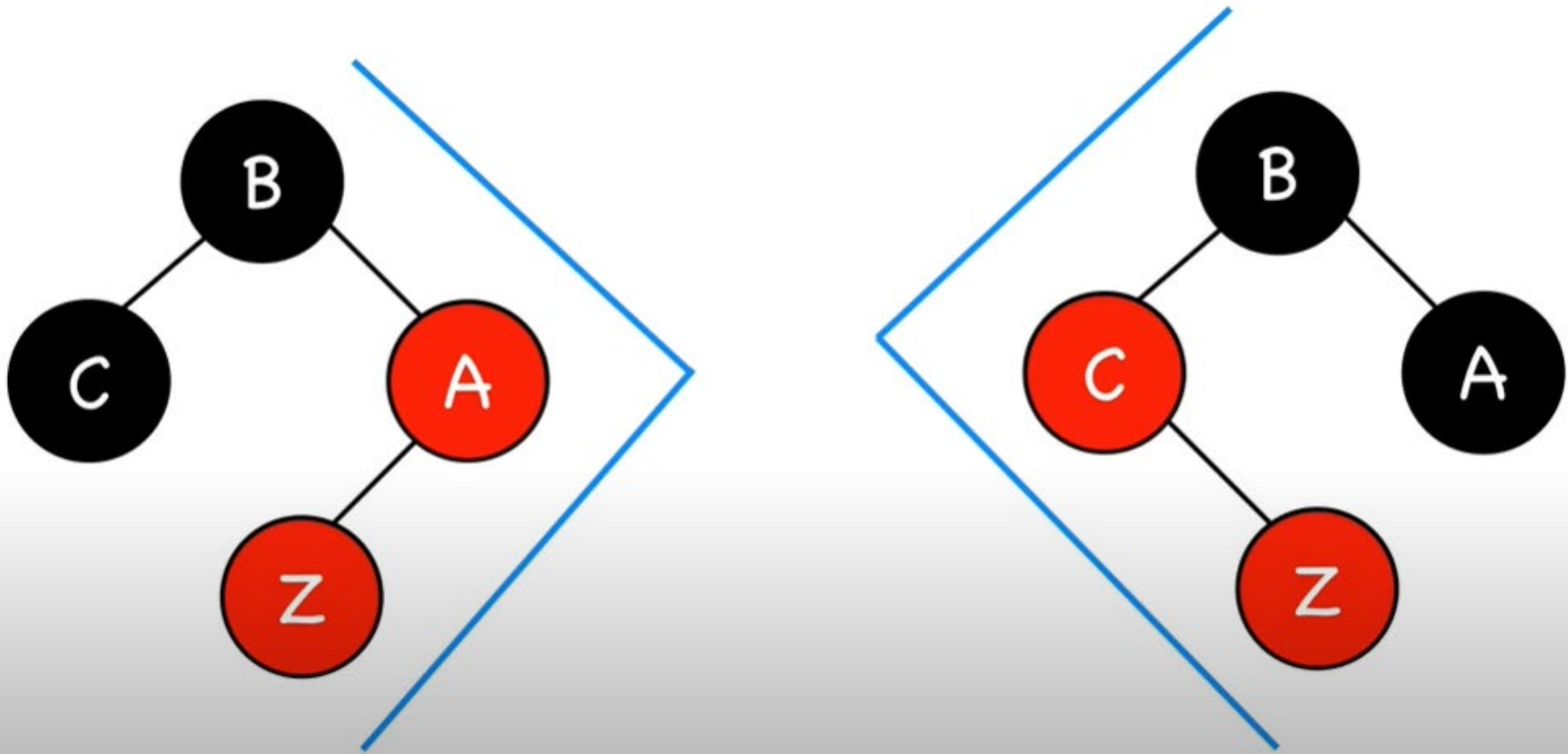
Case 1. Z.uncle = red

- Recolor Z's parents and grandparent



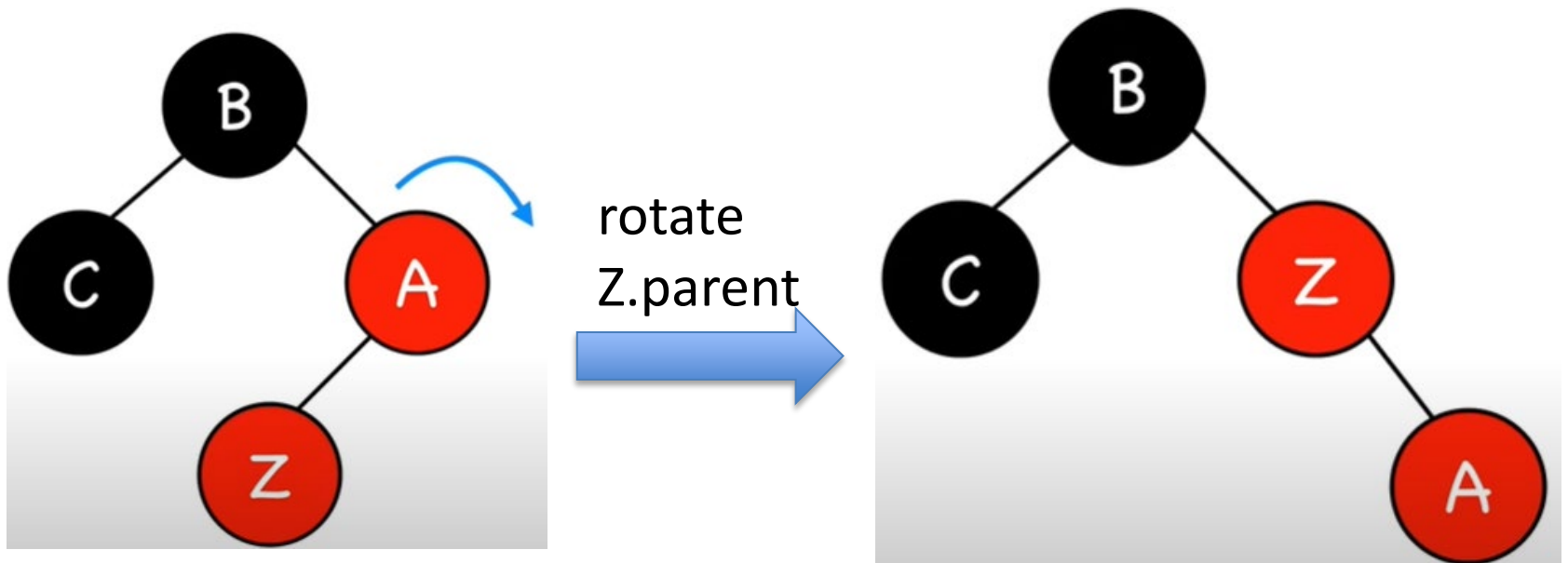
Case 2. Z.uncle = black (triangle)

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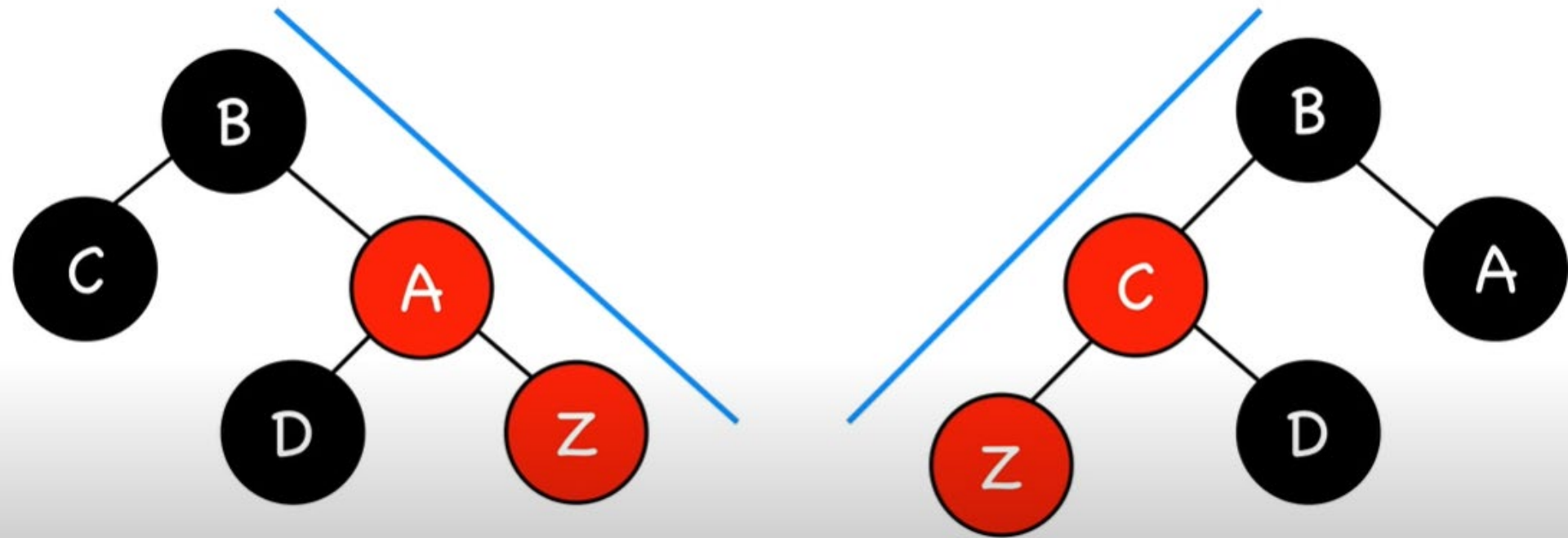
Case 2. Z.uncle = black (triangle)

- Rotate Z.parent
- Turns into Case 3



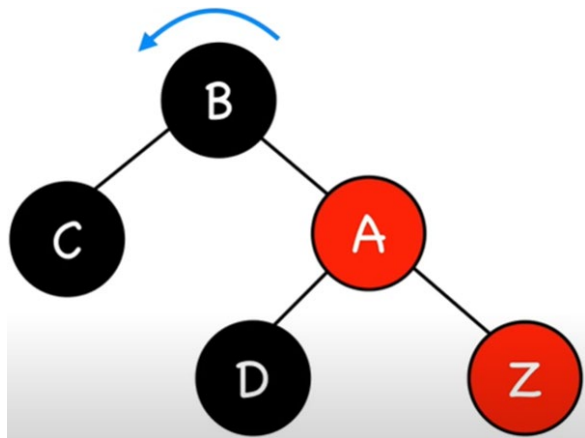
Case 3 Z.uncle = black (line)

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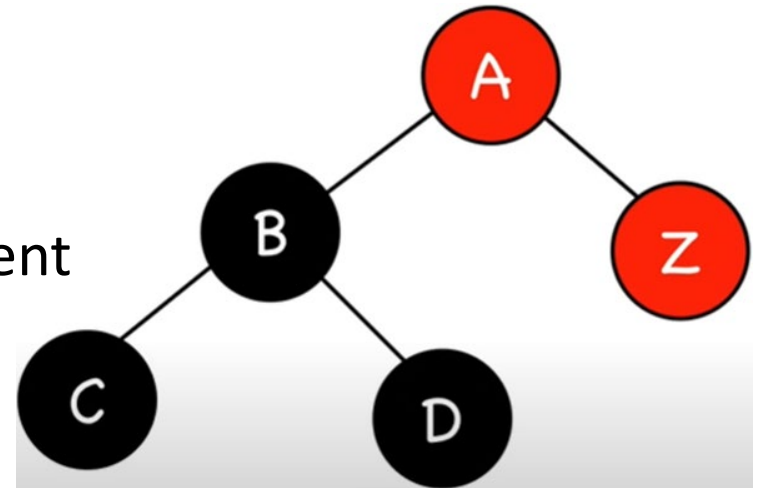


Case 3 Z.uncle = black (line)

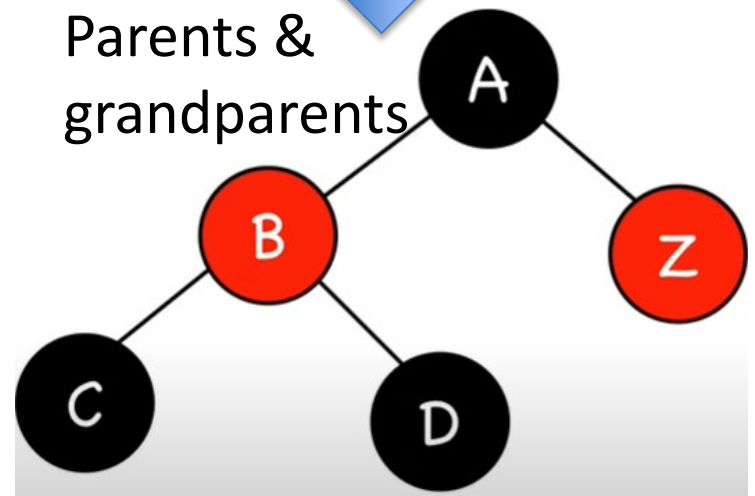
- Rotate Z.grandparent



rotate
Z.grandparent

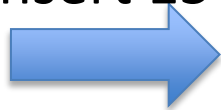


recolor
Parents &
grandparents

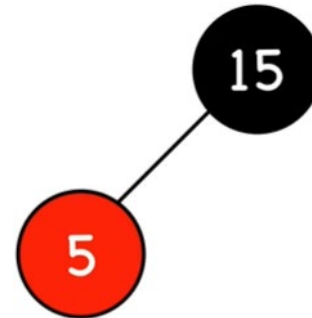
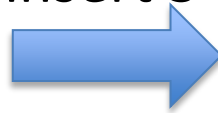


Example 1

insert 15

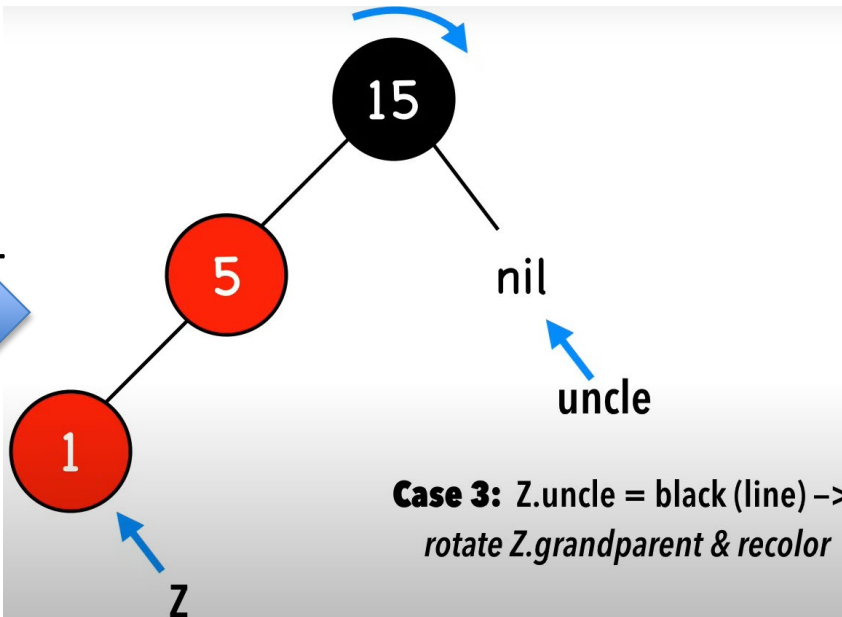
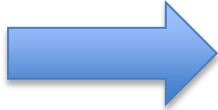


insert 5

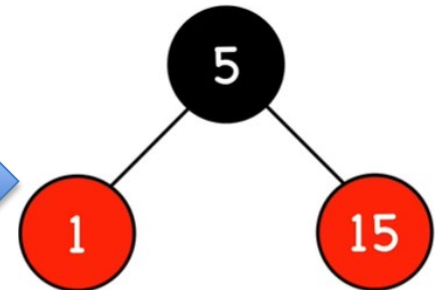
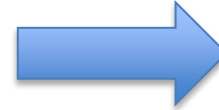


Case 0: $Z = \text{root} \rightarrow \text{color black}$

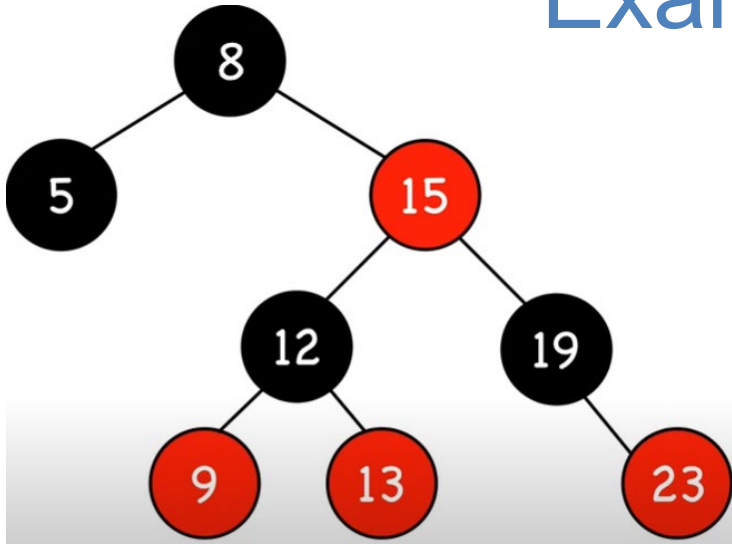
insert 1



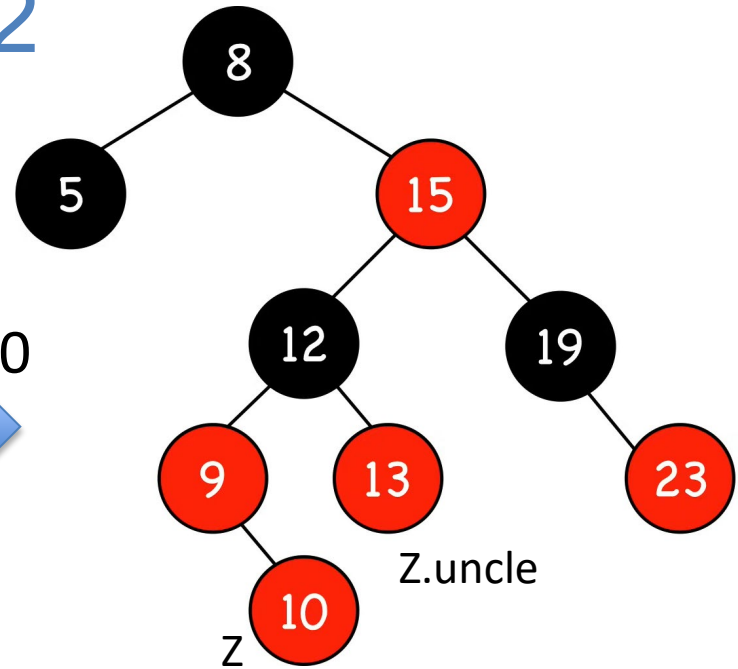
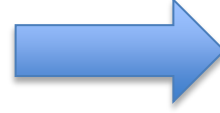
Case 3: $Z.\text{uncle} = \text{black (line)} \rightarrow$
rotate $Z.\text{grandparent}$ & recolor



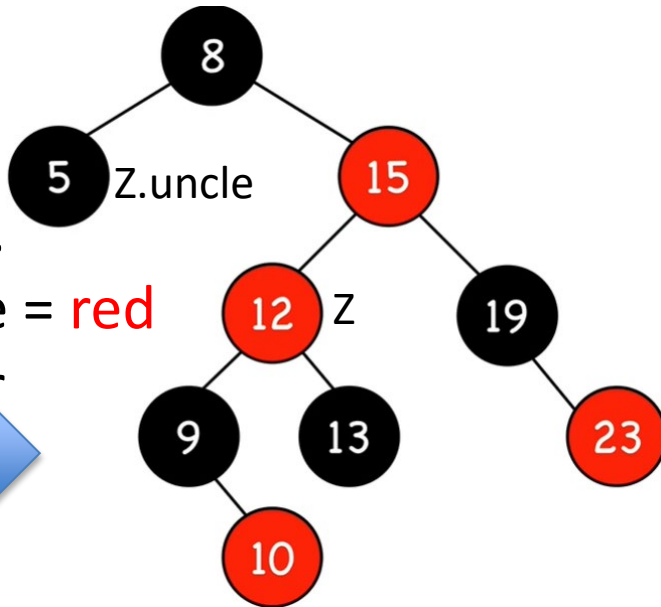
Example 2



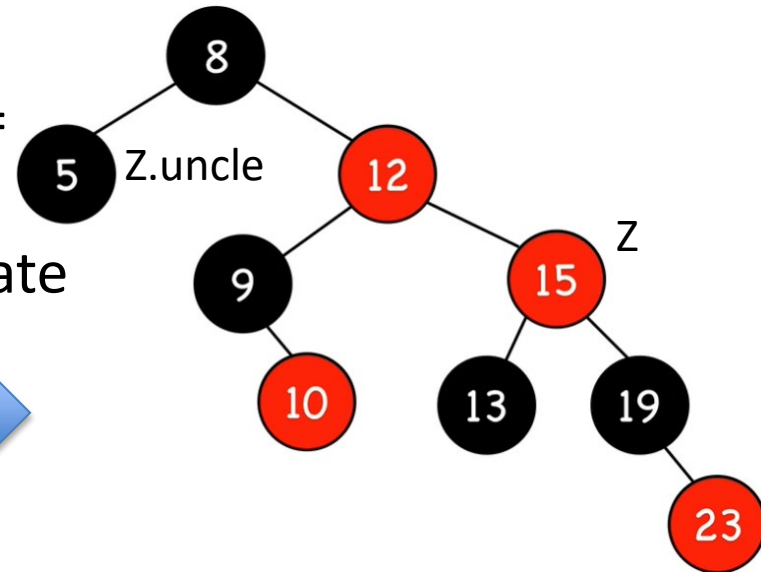
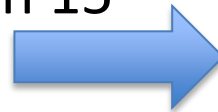
insert 10



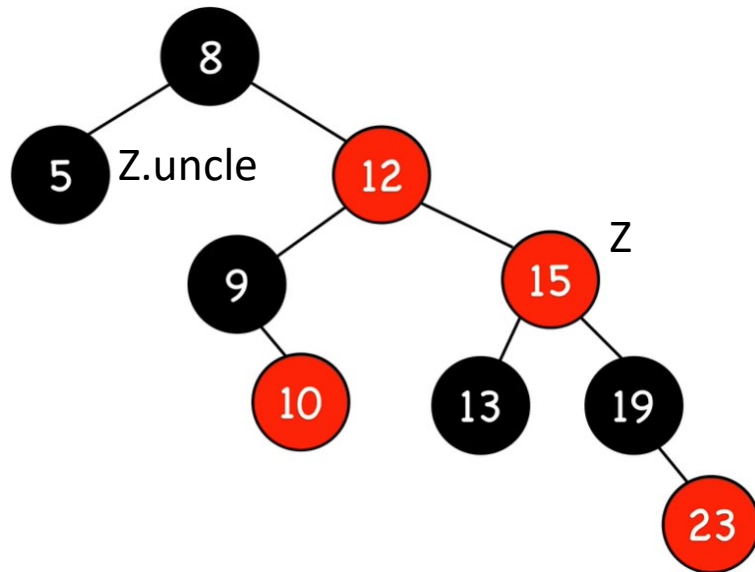
Case 1.
Z.uncle = red
recolor



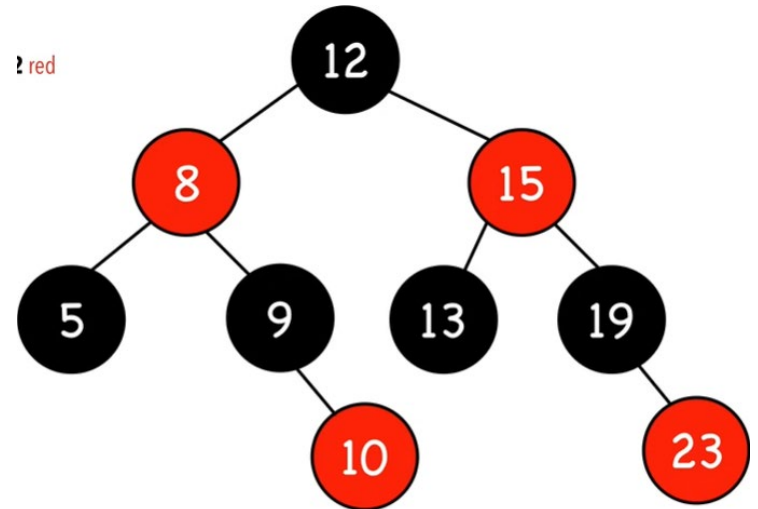
Case 2.
Z.uncle = black
right rotate
on 15



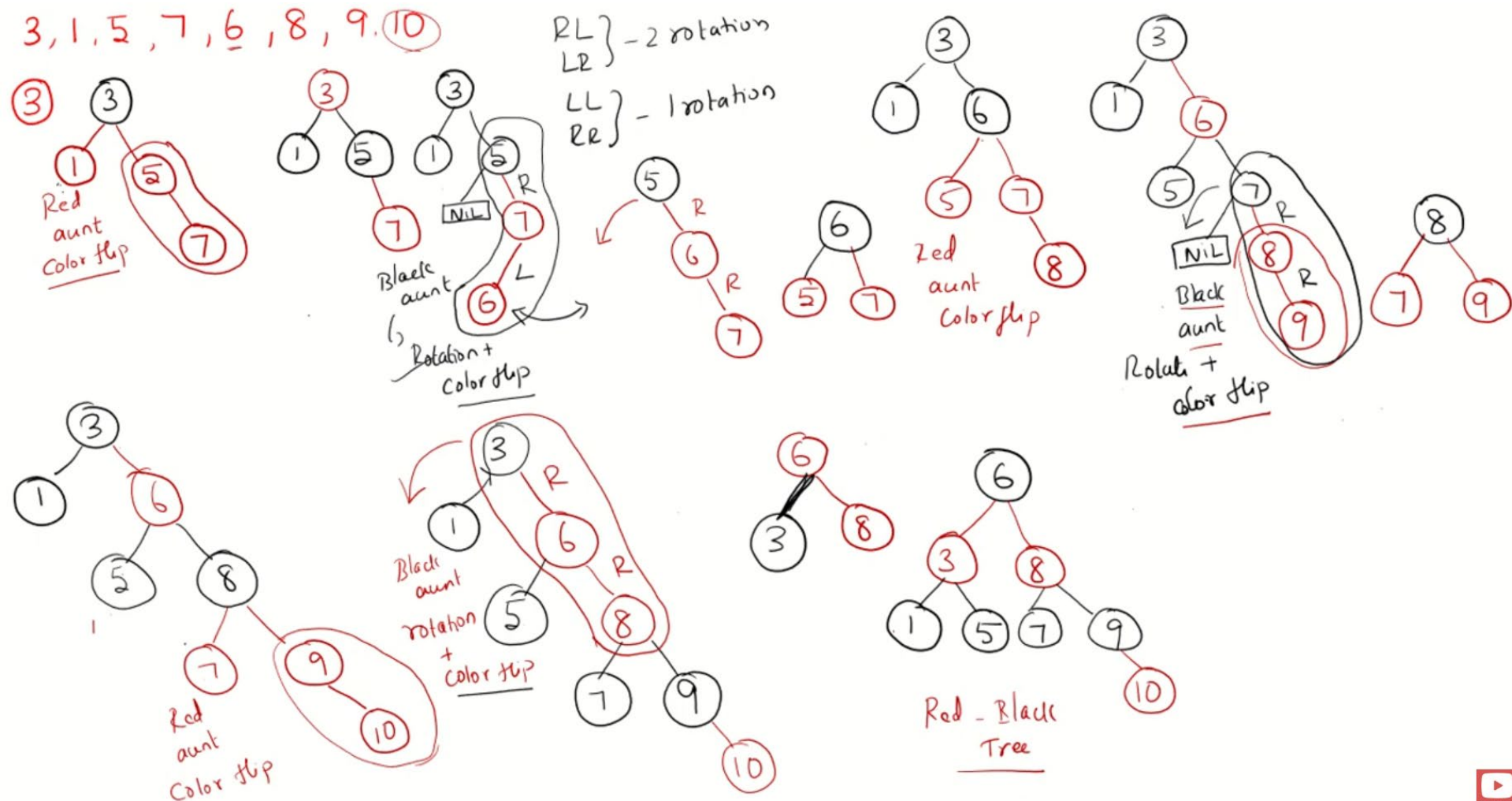
Example 2 Con't



Case 2.
Z.uncle =
black
left rotate
on 8 &
recolor



Another Example



Red Black Tree – Insertion

<https://www.youtube.com/watch?v=9ubIKipLpRU>

Time Complexity

- 1. Insert : $O(\log(n))$
 - maximum height of red-black trees
- 2. Color red : $O(1)$
- 3. Fix violations :
 - Constant # of:
 - a. Recolor : $O(1)$
 - b. Rotation: $O(1)$
- Overall time complexity: $O(\log(n))$

Applications

- Red–black trees are widely used as system symbol tables.
 - Java: `java.util.TreeMap`, `java.util.TreeSet`.
 - C++ STL: `map`, `multimap`, `multiset`.
 - Linux kernel: completely fair scheduler, `linux/rbtree.h`.
 - Emacs: conservative stack scanning.

Video Tutorials

- Red-Black Trees // Michael Sambol
 - https://www.youtube.com/playlist?list=PL9xmBV_5YoZNqDI8qfOZgz_bqahCUMUEin
 - Lecture slides based in this video series
- Red Black Tree – Insertion
 - <https://www.youtube.com/watch?v=9ubIKipLpRU>
- Introduction to Red-Black Tree
 - <https://www.geeksforgeeks.org/introduction-to-red-black-tree/>